# Groebner bases under composition and multivariate matrix factorization 

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## Contents of the talk

In this talk, I will introduce some of my work related to Groebner bases.
\& I.Behavior of a Groebner basis under composition - New Results;
$\diamond$ II. Multivariate polynomial matrix factorization - New progress.

PART(I)
Groebner bases under composition

## Groebner bases under composition

Hoon Hong initiated the study of Greobner bases under composition

## Groebner bases under composition

Early papers:
(i). H.Hong, Groebner bases under composition II, ISSAC, 2006.

- (ii) H.Hong, Groebner basis under composition I, JSC (1998), 25, 647-662.
( (iii) Gutierrez , J. Miguel, R, Reduced Groebner basis under composition, JSC(1998), 26, 433-444.


## Basic research Problem

Let us consider the polynomial ring in $n$ variables: $K[x]=k\left[x \_1, \ldots x \_n\right]$, and a given term ordering >.

The basic problem can be described as follows:

For an endomorphism $U$ of $k[x]$, and a finite subset $G$ of $k[x]$, how can we compute a Groebner basis of $U(G)$ under the term oredering > by means of a Groebner basis of $G$ and $U$ ?

## Hong's Theorem

For any Grobner basis G of k[X], $\mathrm{U}(\mathrm{G})$ is a Groebner basis if and only if (a) for any terms $p$ and $q, p<q$ implies $\operatorname{lt}(\mathrm{U}(\mathrm{p}))<\operatorname{lt}(\mathrm{U}(\mathrm{q}))$;
*(b) It(U) is a permuted powering, that is, every component of $U$ has a power of one variable as the leading term, where $\mathrm{U}=(\mathrm{u} 1, \ldots \mathrm{un}): \mathrm{xj} \rightarrow \mathrm{uj}$. and $\operatorname{lt}(\mathrm{U})=(\operatorname{lt}(\mathrm{u}), \ldots, \operatorname{lt}(\mathrm{un}))$.

## Result of Gutirrez etc.

Gutirrez etc. have proved a result on reduced case. Their result is:
For any reduced Groebner basis $G, U(G)$ is a reduced Groebner basis if and only if
$\diamond(a)$ for any terms $p, q, p<q$ implies $\operatorname{lt}(\mathrm{U}(\mathrm{p}))<\operatorname{lt}(\mathrm{U}(\mathrm{q}))$.
$\diamond(\mathrm{b})$ every component of $U$ is a polynomial in one variable, and different polynomial involves in different variable.

## A generalization

First, We establish a more general framework:

Let $U$ be a homomorphism from $k[x]=$ $K[x 1, \ldots x n]$ to $k[y]=k[y 1, \ldots y m]$ s.t. $n<m$ or $\mathrm{n}=\mathrm{m}$.

Given two term orderings $>_{1}$ on $\mathrm{k}[\mathrm{x}]$ and $>_{2}$ on $\mathrm{K}[\mathrm{y}]$. In this case $u j$ is in $k[y]$.

## A generalization

Under this new framework, Hong's theorem can be generalized as follows:
For any Groebner basis $G$ with respect to $<_{1}, \mathrm{U}(\mathrm{G})$ is a Groebner basis with respect to <2 if and only if
$\diamond(1)$ for any terms $p, q, p<{ }_{1} q$ implies that $\operatorname{lt}(\mathrm{U}(\mathrm{p}))<2 \operatorname{lt}((\mathrm{U}(\mathrm{q}))$;
(2) $\operatorname{It}(\mathrm{ui})$ and $\operatorname{It}(\mathrm{uj})$ are pairwise coprime for different i and j .

## A generalization

But in reduced case, it seems to lack a sufficient and necessary condition in this generalization case.

Details can be found in:
Remarks on Groebner basis for ideals under composition, ISSAC, 2001.

## Check the term ordering conditions

From the above theorems, we need to solve the term ordering compatible problem in order to apply these theorems:

Find an efficient algorithm to check if for any terms $p$ and $q, p<1 q$ implies $\operatorname{lt}(\mathrm{U}(\mathrm{p}))<2$ $\mathrm{lt}(\mathrm{U}(\mathrm{q}))$ ?(Hong)

## Continue

Let T be the matrix corresponding to exponent vectors of leading terms of $u j$, where $\mathrm{U}=(\mathrm{u} 1, \ldots \mathrm{un})$
Let term ordering $>_{1}$ be represented by a nxn matrix $A$, and $>_{2}$ be represented by matrix $B$. Thus our method is using rational elementary transformation to A and TB simultaneously to obtain some standard form in some sense, then check if these standard form are the "same" which is up to a positive number multiple.

## Continue

So called elementary rational transformation for the real matrices we mean:
$\diamond$ I. Multiplying a row of a matrix by a non-zero rational number;

* II. Interchanging any two rows;
* III. Adding a rational multiple of one row to another row

See JSC(2003), vol.35.

## Homogeneous case

An interesting problem is:

Under what conditions that for any homogeneous groebner bases $G, U(G)$ is homogeneous greobner bases?

Journal of Algebra, computational algebra section, In press.

## Homogeneous case

We have provided a complete answer:

For any homogeneous Greobner bases G, U(G) is a homogeneous Groebner bases if and only if
$\diamond(I)$ for any terms p,q with $\operatorname{deg}(p)=\operatorname{deg}(q)$, $\mathrm{p}<\mathrm{q}$ implies $\mathrm{It}(\mathrm{u}(\mathrm{p}))<\mathrm{It}(\mathrm{u}(\mathrm{q}))$;
$\diamond$ (ii) $\operatorname{lt}(\mathrm{U})$ is a permuted powering, and every ui has the same degree.

## Further results

The problems can be further extended more general case. Let $L$ be an arbitrary grading on $\mathrm{k}[\mathrm{x}]$, we may ask:

Under what conditions that for any Lhomogeneous groebner bases $G, U(G)$ is Groebner bases?

## An answer

For any L- homogeneous $\mathrm{Gb} \mathrm{G}, \mathrm{U}(\mathrm{G})$ is GB if and only if
*(I) for any terms $p>q, L(p)=L(q)$ implies $\mathrm{It}(\mathrm{U}(\mathrm{p}))>\operatorname{lt}(\mathrm{U}(\mathrm{q}))$;
$\rangle(\mathrm{II}) \mathrm{It}(\mathrm{U})$ is a permuted powering;

## Remarks

This new result unifies all the previous results. Because if $\mathrm{L}(\mathrm{xi})=0$ for every I, we get Hong's theorem; if $L(x i)=1$ for every I, we get the usual homogeneous result.

The proof is slightly hard; submitted to Journal of Algebra, CA section.

## Summary

$\diamond$ From above descriptions, we see that only universal cases are considered.
$\diamond$ For the basic problem: Given a Groebner basis G, and a homomorphism U, How can we compute the Groebner basis of $\mathrm{U}(\mathrm{G})$ by means of $G$ and $U$ ?
$\diamond$ We have not solved the basic problem! It seems to have some difficulties.

## Summary

- There are a lot of related work in

Resultant, subresultant, Sagbi bases, non-commutative Groebner bases etc.
-See related papers.

PART(II)

Multivariate polynomial matrix factorization

## Basic Problem

Given a nD polynomial matrix $F$ with full row rank of size (I, m).

Let $d$ be the gcd of all the I-minors, f|d, whether or not there exists square matrix G and matrix F1,such that

$$
\mathrm{F}=\mathrm{GxF} 1 \text { with } \operatorname{det}(\mathrm{G})=f \text { ? }
$$

## The History

$\diamond$ When $\mathrm{n}=1$, this is solved completely using Gauss elimination, and play an essential role in linear systems and control theory, convolutional codes.
$\diamond$ When $\mathrm{n}=2$, this is solved using pseudodivison by N. K. Bose and others, used in 2-D systems and signal processing.
$\geqslant$ When $n>=3$, this is thought as very hard! Because previous methods can not be generalized to this case.

## Motivation for research

Why we concern about this problem?
$\diamond$ (i). Generalize linear system theory to multidimensional systems theory since 1976.
$\diamond$ (ii). Many problems in multidimensional systems and signal processing can be formulated in multivariate matrix problems.
$\diamond$ (iii). May be useful in multidimensional convolutional codes.

* (iv). Possibly other applications.


## Reduced minors

In order to introduce Lin-Bose Problem, we give a basic definition as follows:

Let F be a full row rank matrix of size (l, m), let a1,...,ak be all the Iminors of $F, d$ be the $g . c . d$ of $a 1, \ldots, a k$, and let bj such that $a j=d x b j$. b1,...,bk is said to be the reduced minors of $F$.

## Some research papers

There are many papers concerning this problems, we just mention:
$\diamond$ C. Charoenlarpnopparut, N. K.Bose, Multidimensional FIR filter Bank Design using Groebner basis, IEEE Trans. Circuits Systems II Vol 46, 1999, pp1475-1486.
$\diamond$ Z. Lin, Notes on nD polynomial matrix factorization, Multidimensional system and Signal Processing, Vol. 10, 1999, 379-393.

## Research papers

ZZ. Lin, Further results on nD polynomial matrix factorization, Multidimensional system and system processing, Vol 12, 2001, pp199-208.

- M. Wang, C.P.Kwong, Computing GCLF using Syzygy algorithm, Proceedings of ICMS, 2002.


## Research papers

- M.Wang,D.Feng, On Lin-Bose problem, Linear algebra and application, vol. 290, 2004;
*M. Wang, C.P.Kwong, On multivariate matrix factorization problems. Mathematics of Control, signals and systems. 2005.


## Zero prime factorization

A full row rank matrix is zero prime if all its maximal order minors generate the unit ideal.

Lin-Bose propose the following problem:
Given F, full row rank of size IXm, if all the reduced minors generate the unit ideal, then whether or not $F$ can be factorized as $F=G x F 1$, with $\operatorname{det}(G)=d$, $d$ the $g c d$, and $F 1$ being zero prime?

## Zero prime factorization

J.F.Pommaret gave a proof using algebraic analysis method in 2001 Euro control conference, his method only holds for the complex number field.

We give a full proof for any field. Linear algebra and applications, 2004

## Minor prime factorization

Let $F$ be as above, $d$ the gcd of all the Iminors of $F$. Whether or not there exsits $G$ such that $F=G x F 1$ and $\operatorname{det}(G)=d$ ?
If such factorization exists, we call it minor prime factorization of $F$.
Several people gave counter-examples to showed that the above problem may have no solution in 1976-1979;
However, no one propose a sufficient and necessary condition!

## Minor prime factorization

We first found a sufficient and necessary condition.

Let K be submodule generated by all the rows of F . There exists a minor prime factorization if and only if the colon submodule K:d is a free module of rank I.

## Remarks to the proof

The proof of the above result relies on a characterization of so-called minor left prime(MLP) matrix.
A full row rank matrix $F$ of size $1 x m$ is said to be a MLP matrix if all the I-minors have only trivial common divisors.

The result just holds for any field. See Mathematics of control, signals and systems, 2005 for details.

## Remarks to proof

We first prove, a matrix $F$ is a MLP if and only if $\mathrm{K}: \mathrm{d}=\mathrm{K}$, where $\mathrm{K}=$ submodule generated by all the rows of $F, d$ the g.c.d of all the $|x|$ minors of $F$.

Then using a lifting to the linear mapping, we get the factor.

Note that when F is not full rank, this result will not hold, and above result does not hold either.

## An algorithm

Based on above theorem, we propos an algorithm as follows:
Step 1 Let F as above, d the g.c.d of all I times I minors of $F$. Compute the syzygy module of row vectors of $F$ and -dxI_m. If this syzygy module is free of rank I, then form a generating matrix $[\mathrm{G} \mid \mathrm{H}]$ such that G is a $\mid \mathrm{xl}$ matrix, thus $\mathrm{dG}^{\wedge}\{-1\}$ is the desired matrix factor;
Step 2 If this syzygy module is not free, then there do not exist such a matrix factor.

## Continue

We can easily check if a submodule is free using Fitting ideals.

## Factor prime matrix

$F$ is said to be factor prime, if for any square matrix $G, F=G x F 1$, implies that $\operatorname{det}(\mathrm{G})$ is a non-zero constant.

A basic problem is to check if $F$ is factor prime matrix .

## Factor prime matrix

Checking if a matrix is factor prime is a longstanding open problem since the concept was proposed in 1979.
D.C.Youla, G,Gnavi, Notes on n-Dimensional systems theory, IEEE circuits and systems, vol. CAS-26, 1979.

We have also obtained some partial answers!

## Remarks for checking of FLP

No any other methods available to attack deciding problem of factor prime matrix;
This is the first result concerning factor prime matrix.
Full solutions can not obtained at this time, to be submit.

## Remarks to the algorithm

Major problem is : how to extract a system of generators with I elements from a system of generators of syzygy module?

Book "Computational Commutative algebra I"by Martin Kreuzer and Lorenzo Robbiano provides a partial answer in Corollary 3.1.12 which was implemented in CoCoA.

## Compare with other related work

*Previous papers only deal with very special cases.
*Above results are true for any field, that is, without restriction to characteristic zero and algebraically closed.
*Only Groebner bases theory is available to attack matrix factorization problems currently.

## Some problems

I'd like to list some problems which deserve to be considered :
Problem 1: if there are other efficient methods to deal with this problem?

- Problem 2: How to find a minimal generating system from a given generating system?
- Problem 3: If there are other criteria to assure that the existence of a generating system of I elements? (I=rank)


## Some problems

Problem 4: If elementary transformations for multivariate polynomial matrices can be used to find matrix fator?

- Problem 5: Find other characterizations for the existence of the matrix factor!
Problem 6: Find the methods for dealing with non-full rank cases.


## Conclusion

$\diamond$ Matrix factorization problems connects with many problems in system theory, for example, rational matrix fraction description.
$\diamond$ Generalization of convolutional codes in one variable to multidimension also needs to develop multivariate matrix theory.

- It needs to develop new algorithmic methods!


## Thank you for your attention!

